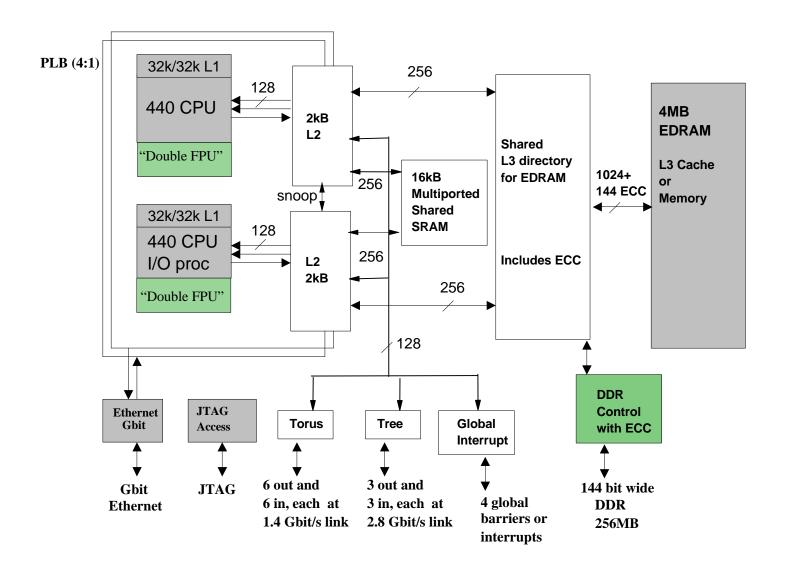


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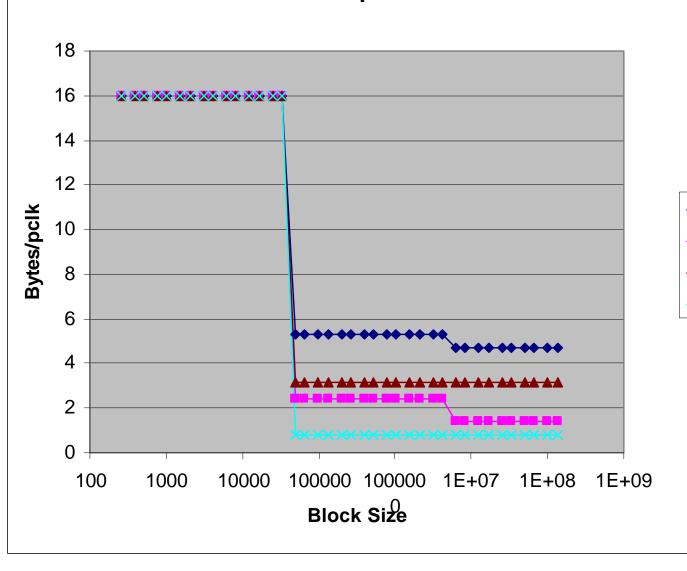


## Some Features

- no MESI, but
- hardware support for
  - flushing L1 cache
  - locks/barriers
  - fast core-to-core signaling (interrupt)
- 16kB SRAM + L3 scratchpad
- fine control over memory access behavior on all memory hierarchy levels



### **Sequential Read Bandwidth**



→ L2 enabled, L3 enabled

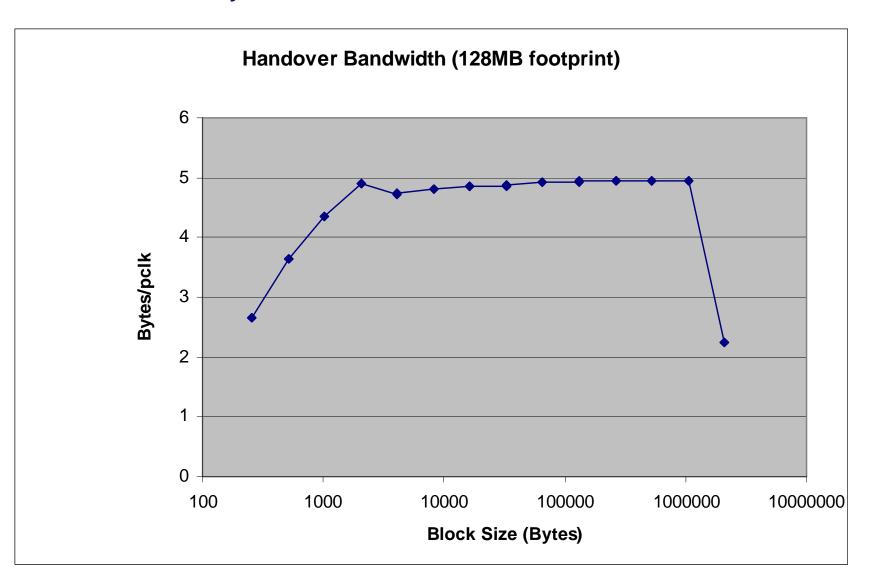
L2 disabled, L3 enabled

L2 enabled, L3 disabled

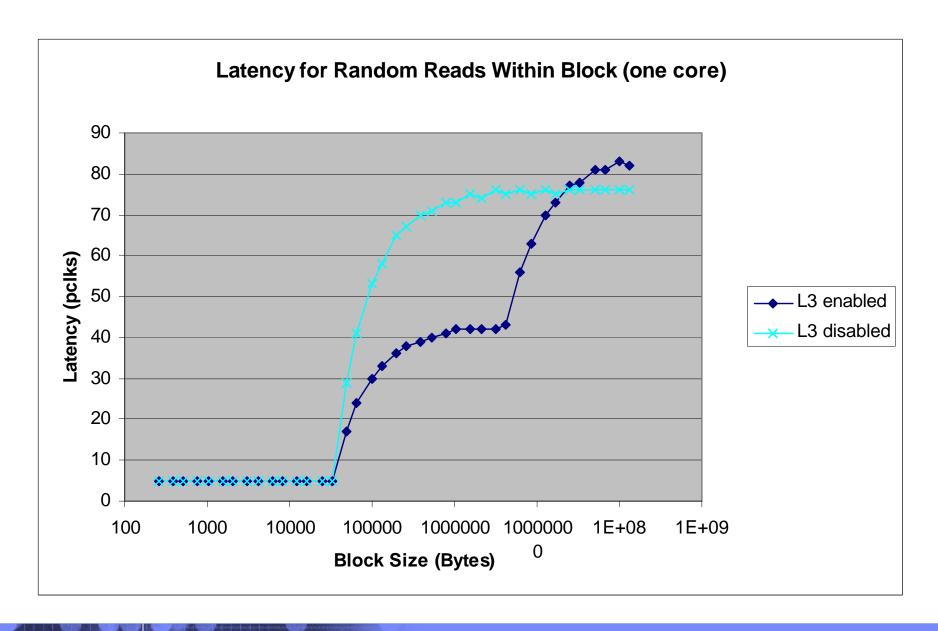
L2 disabled, L3 disabled



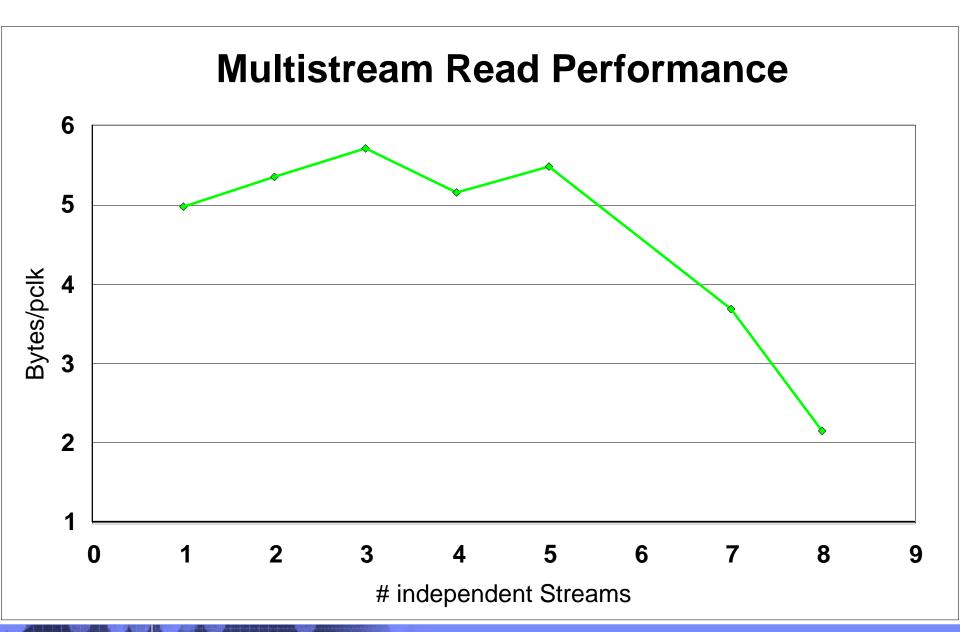
#### Delivered Memory Bandwidth Between CPUs Sustains Maximum Link Bandwidth



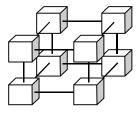






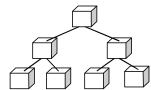






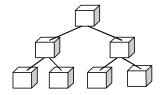
#### **3 Dimensional Torus**

Point-to-point



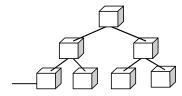
#### **Global Tree**

Global Operations



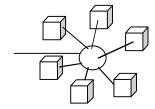
#### **Global Barriers and Interrupts**

• Low Latency Barriers and Interrupts



#### **Gbit Ethernet**

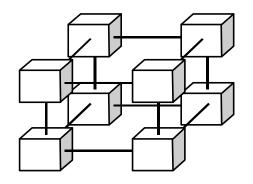
• File I/O and Host Interface

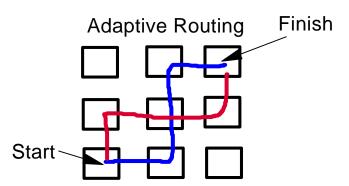


#### **Control Network**

Boot, Monitoring and Diagnostics







- 32x32x64 connectivity
- Backbone for one-to-one and one-to-some communications
- 1.4 Gb/s bi-directional bandwidth in all 6 directions (Total 2.1 GB/s/node)
- 64k \* 6 \* 1.4Gb/s = 68 TB/s total torus bandwidth
- 4 \* 32 \*32 \* 1.4Gb/s = 5.6 Tb/s Bisectional Bandwidth
- Worst case hardware latency through node ~ 69nsec
- Virtual cut-through routing with multipacket buffering on collision
  - Minimal
  - Adaptive
  - Deadlock Free
- Class Routing Capability (Deadlock-free Hardware Multicast)
  - Packets can be deposited along route to specified destination.
  - Allows for efficient one to many in some instances
- Active messages allows for fast transposes as required in FFTs.
- Independent on-chip network interfaces enable concurrent access.



### Link Utilization on FFT Communication Pattern is above 97%

### All-to-All Efficiency for Mesh & Torus Topologies for BlueGene/L

All messages are 100 packets, each packet 256B

	32 (4x4x2)	512(8x8x8)
Mesh Time (Processor cycles)	3.9M	149.8M
% Peak	88%	74%
Torus Time (Processor cycles)	1.8M	56.2M
% Peak	96.8%	98.3%
Torus Theoretical Bound without idle packets	1.7M	55.3M



#### Broadcast in a 2-D plane achieves 98.8% of Peak (4x4 plane)

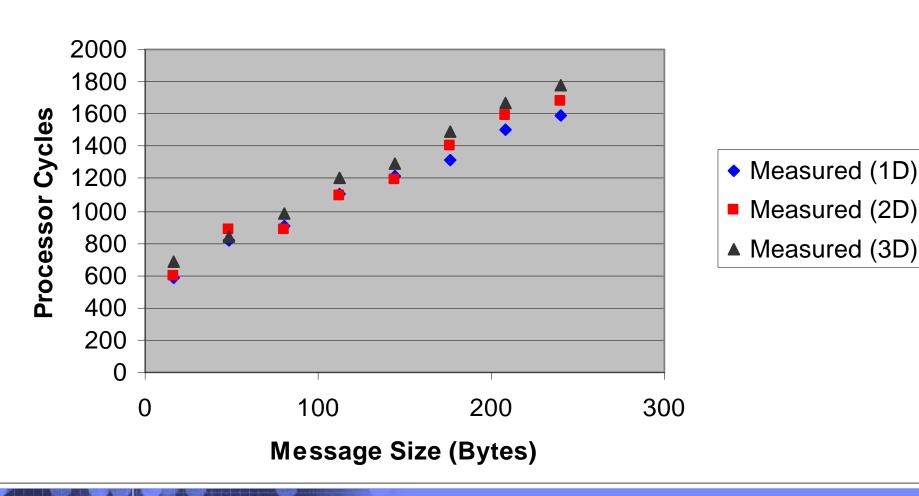
Torus Network supports Hardware Multicast
 Fundamental to dense solvers
 Confirms network efficiency used in Linpack benchmark efficiency model

Number of Packets	Payload (MBytes)	Measured (pclks)	Theoretcial (pclks)	Effciency
10,000	2.4	2,727,674	2,700,000	98.9%
100,000	24	27,234,320	27,000,000	99.1%



### Prototype Delivers ~1usec Ping Pong low-level messaging latency



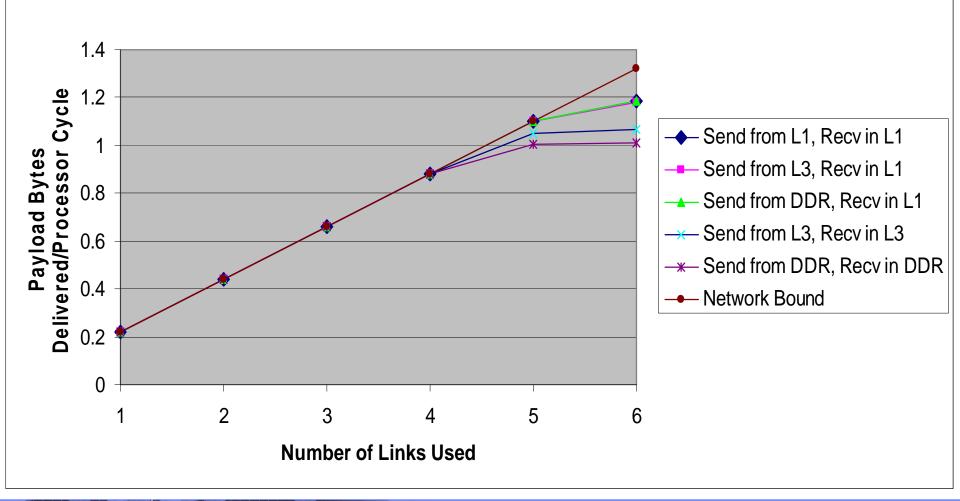




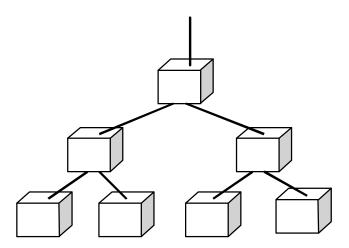
#### Nearest neighbor communication achieves 75-80% of peak

## **Torus Nearest Neighbor Bandwidth**

(Core 0 Sends, Core 1 Receives, Medium Optimization of Packet Functions)





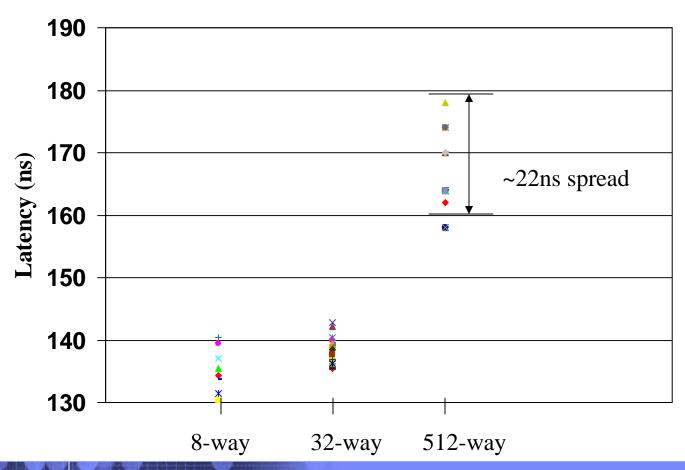


- Four Independent Barrier or Interrupt Channels
  - Independently Configurable as "or" or "and"
- Asynchronous Propagation
  - Halt operation quickly (current estimate is 1.3usec worst case round trip)
    - > 3/4 of this delay is time-of-flight.
- Sticky bit operation
  - Allows global barriers with a single channel.
- User Space Accessible
  - System selectable
- Partitions along same boundaries as Tree, and Torus
  - Each user partition contains it's own set of barrier/ interrupt signals



## Global Barriers: Latency and Spread

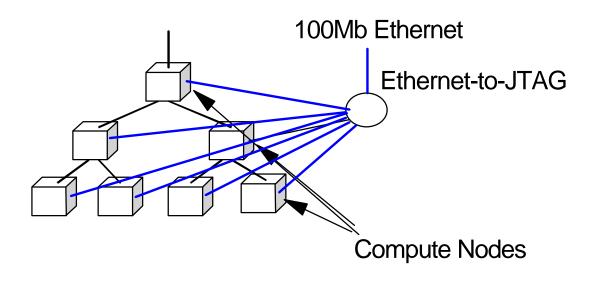
Measured results from prototype confirms ability to scale to 64K with sub-microsecond barrier and interrupt latency



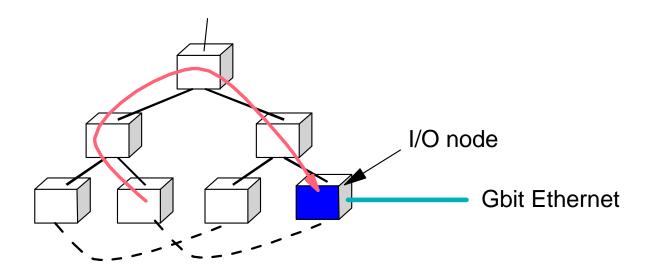


#### **JTAG** interface to 100Mb Ethernet

- direct access to all nodes.
- boot, system debug availability.
- runtime noninvasive RAS support.
- non-invasive access to performance counters
- Direct access to shared SRAM in every node







#### Gb Ethernet on all I/O nodes

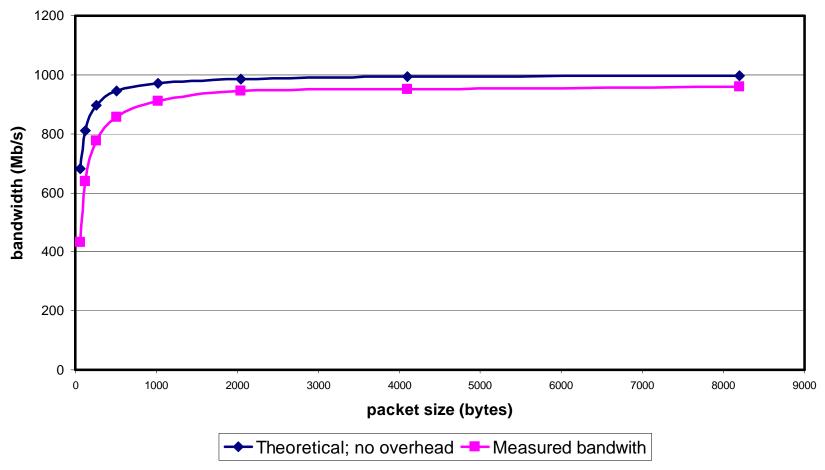
- Gbit Ethernet Integrated in all node ASICs but only used on I/O nodes.
- Funnel via global tree.
- I/O nodes use same ASIC but are dedicated to I/O Tasks.
- I/O nodes can utilize larger memory.

## Dedicated DMA controller for transfer to/from Memory Configurable ratio of Compute to I/O nodes

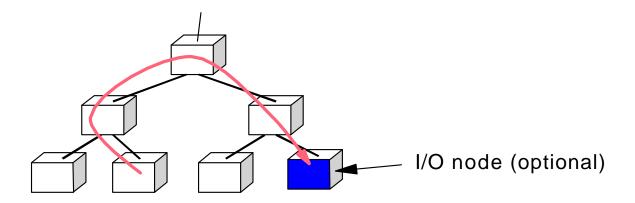
I/O nodes are leaves on the tree network



## Ethernet Performance is at 97% of Peak







- High Bandwidth one-to-all
  - 2.8Gb/s to all 64k nodes 68TB/s aggregate bandwidth
- Arithmetic operations implemented in tree
  - Integer/ Floating Point Maximum/Minimum Integer addition/subtract, bitwise logical operations
- Latency of tree less than 2.5usec to top, additional 2.5usec to broadcast to all
- Global sum over 64k in less than 2.5 usec (to top of tree)
- Used for disk/host funnel in/out of I/O nodes.
- Minimal impact on cabling
- Partitioned with Torus boundaries
- Flexible local routing table
- Used as Point-to-point for File I/O and Host communications



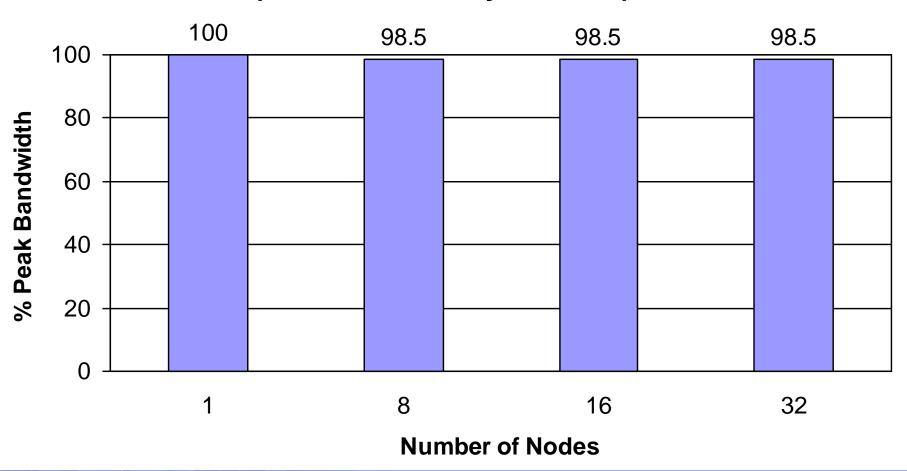
## TREE features cont...

- 350 MB/s per link per direction
   3 network ports -> total: 2.1 GB/s
- Reliable Packet Transport (24bit Packet CRC)
  - + 32bit Link CRC
  - + Injection/Reception Checksums
- Cut-through routing (~60ns latency per hub)
- 256 Bytes payload per packet
- 16 classes / routes per node
- 2 independent virtual networks
- integrated ALU (AND,OR,XOR,MAX,ADD)Datatypes: 16bit ... 2048bit (unsigned integer)
- Global combine operation 64k -> 1 in less than 2.5 us



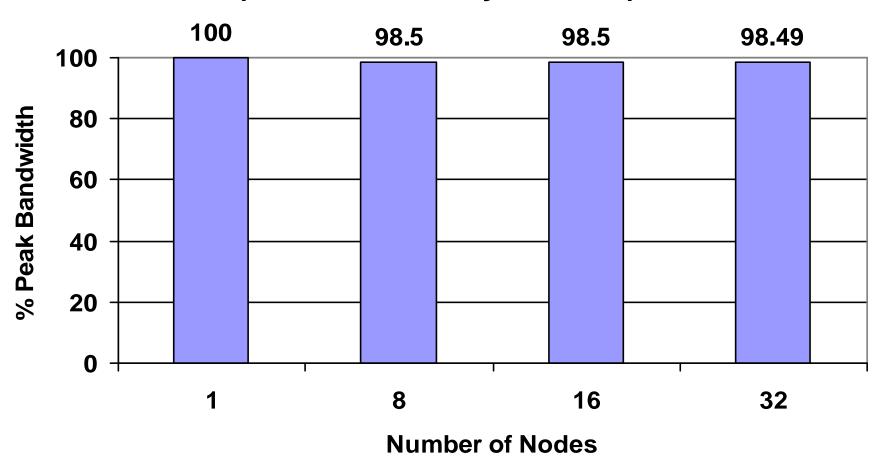
Delivered Broadcast bandwidth on Tree Network is consistent with pipelined architecture and should scale to 64K

## Bandwidth: Streaming Broadcast (Peak = 0.485 Bytes/Pclk)



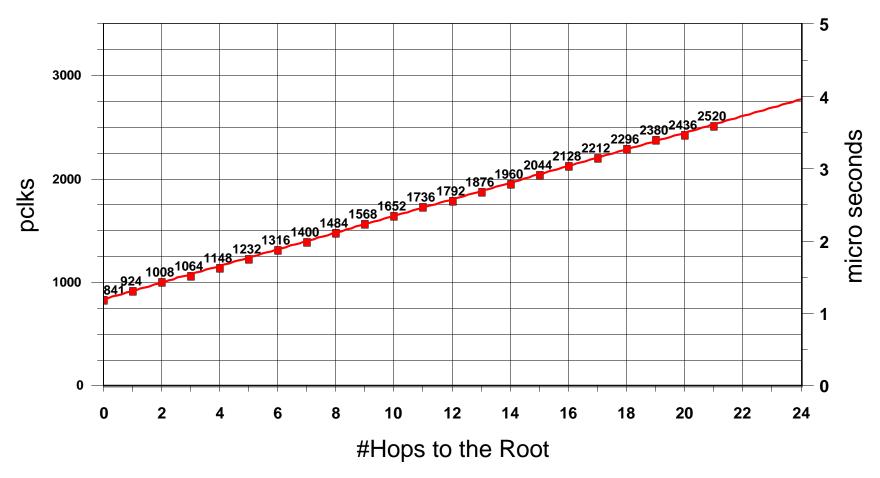


# Bandwidth: Streaming Reduction (Peak = 0.485 Bytes/Pclk)









R-square = 1 # pts = 17  
$$y = 837 + 80.5x$$



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